



Projects Selection In Knapsack Problem By Using Artificial Bee Colony Algorithm

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Abstract

One of the combinatorial optimization problems is Knapsack problem, which aims to maximize the benefit of objects whose weight not exceeding the capacity of knapsack. This paper introduces artificial bee colony algorithm to select a subset of project and represented by knapsack problem to put the best investment plan which achieve the highest profits within a determined costs, this plan is one of the applications of the financial field. The result from the proposed algorithm implemented by matlab (8.3) show the ability to find best solution with precisely and rapidity compared to genetic algorithm.

1- Introduction

The most popular subset selection problem is knapsack problem (KP). The goal is to set a subset of n items from N original items to a knapsack, such that the summation of profit of the selected items is maximized which sum of weight not exceeding the capacity of knapsack [1].

Knapsack problem has many practical applications are not limited to packing problems: suppose that n projects are available to an investor and that the profit obtained from the j^{th} project is p_j , $j=1,2,\dots,n$, whose costs w_j to invest in project j , by solving (0-1) knapsack problem an optimal investment plan may be found[2].

Many methods such as: branch and bound, cutting planes, and dynamic programming etc., and nature inspired algorithm have been proposed to solve the knapsack problem [3,4].

Lately, evolutionary algorithms such as genetic algorithm (GA)[5,6], particle swarm optimization algorithm (PSO)[7], ant colony optimization algorithm (ACO)[8] and artificial bee colony algorithm (ABC)[4] efficiently have been developed

for solving knapsack problem and reaches to the best solution.

In this paper uses ABC algorithm which proved efficiently in select subset of projects in knapsack problem to build best investment plan compared by GA. The rest of paper is organized as follows: section 2 gives related work, section 3 presents description to solve KP by use ABC algorithm, section 4 views the description of the proposed algorithm, experimental result and conclusion are introduced in section 5 and 6.

2- Related works

Many approaches are proposed to solve knapsack problem to obtain the best solution. In [6] introduced the basic form of genetic algorithm to solve knapsack problem. In[4] presents the using of a novel artificial bee colony algorithm to solve knapsack problem, a hybrid probabilistic mutation schemes in this approach are used to guide the search quickly. In [9] Karboga proposed ABC algorithm which its works similar to a bee colony. In [3] introduced a novel binary biogeography-based optimization algorithm

for the knapsack problem, the result from [3] show that the new method is effective and efficient. In [5] implement a fast and efficient genetic algorithm to solve 0-1 knapsack problem feasibility and effectively. In [10] presents a discrete artificial bee colony for multiple knapsack problem, the result from that the presented method has enhanced convergence speed and quality than other evolutionary algorithms. In [11] applied bee algorithm to solve knapsack problem and attempt to find the best solution in knapsack problem by the number of random flight.

3- Description to solve knapsack problem by using ABC algorithm

Karboga in 2005 [9] for real parameter optimization proposed the ABC algorithm, which simulates the forging behavior of a bee colony. The colony of artificial bees consists of three groups: employed, onlookers, and scout bees. Randomly search for food source positions (solutions) by employed bees. Then, they share information about food source such as nectar amounts (quality of solutions) with the onlookers bees which waiting at the dance area at the hive[12-14]. For colony communication the dance is essential and contains three pieces of information regarding of a flower patch: the direction in which it will bees crossing, its distance from the hive and its quality (i.e., fitness). Provided that the amount of nectar of the new selected source is higher than the previous one, the new source position is memories by the bee, and forgets the old one. Otherwise, she keeps in her memory the previous position. The dances of employed bees are watch by onlooker and decide to choose the scout bees. Then employed bees start to search in the neighborhood area of the sources of food found by scout bees [10].

To solve knapsack problem involves packing the knapsack with projects which have a weight capacity C, and the number of projects 1,2,...,n, each project has weight corresponds to weight(wi) and profit (pi), the goal is to find a set of projects whose weight does not exceed the capacity of knapsack and to maximize the overall profit.

Mathematically, the zero-one knapsack problem can be represented by a vector of binary values X1,X2 Xn, Where Xi = 0 or 1 (1 ≤ i ≤ n). given (C > 0), Wi > 0, Pi > 0, 1 ≤ i ≤ n, the zero-one knapsack problem can be represented by a vector of binary values X1,X2 Xn, Where Xi = 0 or 1 (1 ≤ i ≤ n). satisfy the constraint in equation 1 is the goal from finding a vector [5].

$$\sum_{i=1}^n Wi Xi \leq C \quad \dots \dots (1)$$

And maximize the total profit as:

$$\sum_{i=1}^n Pi Xi \quad \dots \dots (2)$$

To select project by using ABC algorithm a possible solution for the problem can be represented by a food source, and the amount of neater for this food source meet the sum of profit for the selected projects in knapsack. Each bee in ABC algorithm is represented as a string of binary value each bit corresponds to

absence or presence of that project in the solution of the bee [4].

One of approach to construct the initial solution by calculate heuristic information by divide profit to weight as in eq. 3 for all item before entered in steps of ABC algorithm.

$$Ri = Pi/Wi \quad \dots \dots (3)$$

And compute the probability of presence each item in initial solution from eq. 4.

$$Ppi = \left(\frac{C}{T}\right) \cdot \left(\frac{Ri}{Mean(R)}\right) \quad \dots \dots (4)$$

Where the sum of weight for all original items corresponds to T and C represents capacity of the knapsack and Mean(R) is the mean of the relation computed from eq. 4.

After finding a food source by each bee in each iteration, it must be modified where the sum of weight for overall items not exceeds the capacity of knapsack, then evaluate the gathered solution generated by employed bee by compute the fitness (nectar amount) of the kth bee corresponds to the sum of the total profit of the selected items in knapsack by that bee, then apply the selection strategy to choose the best bee as scouts the chance of selecting a bee by onlooker computed from eq.5 [10].

$$Chk = Fk / \sum_{i=1}^M Fi \quad \dots \dots (5)$$

F_k: is the profit of knapsack represented by the kth bee.

$\sum_{i=1}^N Fi$: is the total amount of achieved nectar by all the explorer bees.

M: number of employed bee.

Then continue in running ABC algorithm which includes the following algorithm[15][4]:

1- Initialization: parameters of ABC algorithm are initialized K, N, CC_i =0 (i=1,...,K).

2- Initial solution: generate initial solution (food sources) { FS_i(0) | i=1,...,K } by calculate the value of profit to weight by apply eq.3, Then compute probability of presence of each item in initial solution from eq.4.

3- Loop:

For r=1:N do

{

(1) For i=1:K do (select food sources by employed bees and run local search respectively).

{

- Assassinate each employed bee with a food sources FS_i(r) and calculate its amount of nectar.

- In the neighborhood of FS_i(r) find new FS_i'(r) and compute its amount of nectar;

- Put the better one in { FS_i'(r), FS_i(r) } as a new position of the employed bees;

}

(2) For j=1:K do (perform further local search by onlooker bees which help employed bees)

{

- A food source FS_i(r)selected from { FS_i(r) } for every onlooker bee;

- A new food source FS_i'(r) in the neighborhood of FS_i(r) is find then calculate its nectar amount;

- Consider the better one in $\{ FS_i'(r), FS_i(r) \}$ as a new position of the corresponding bees;
 }
 (3) A new food sources exploiting (scout bees randomly perform global search).
 For $i=1:K$ do (food sources)
 {
 If $FS_i(r) = \{ FS_i(r-1) \}$ then $CC_i = CC_i + 1$;
 If $CC_i = \text{Limit}$ (Limit= $K * N$) then
 {
 - The $FS_i(r)$ is abandon and the associated employed bee become scout;
 - Generate randomly a new $FS_i(r)$ and the scout becomes an employed bee again ;
 - $CC_i = 0$;
 }
 }
 - Register the best food source FS_{best} found so far; $r=r+1$;
 }
 4- output: if the predefined end condition are met then return FS_{best} .

4- Description of the proposed algorithm

In this paper, ABC algorithm used to select subset of projects in KP to build good investment plan, these subset represent the best solution which have maximized gross-profit and the sum of total cost not exceeds the capacity of knapsack. There are set of projects available to investor each project has value of sale-revenue and value of cost-making which consist from four types: cost of direct materials, the cost of direct wages, variable industrial costs and non-variable industrial costs, the goal from proposed algorithm select best subset of projects from this set in less executing time.

The proposed algorithm contains the following steps
 1- Initialization: parameters of ABC algorithm are initialized K =number of population=no. employed bee=no. onlooker bee, N =maximum iteration, CC_i =Abandonment counter=0 ($i=1, \dots, K$).
 2- Enter the array of sales-revenue which size [1...no. of projects], and matrix of cost-making whose dimension [no. of project, 4 = types of cost-making].
 3- Calculate the total cost from summation of cost-making for each project.
 4- Compute Gross-profit for each project by subtract total cost from sales-revenue.
 5- Initial solution: generate initial solution (food sources) $\{ FS_i(0) \mid i=1, \dots, K \}$ by calculate the value of profit to weight by apply eq.3, Then compute probability of presence of each item in initial solution from eq.4, and generate initial solution based on the computed value of probability.

6- Loop:
 For $r=1:N$ do
 {
 (4) For $i=1:K$ do (select food sources by employed bees and run local search respectively).
 {

- Assassinate each employed bee with a food sources $FS_i(r)$ and calculate its amount of nectar (corresponds to the sum of the gross-profit of the selected project in knapsack by that bee).
 - In the neighborhood of $FS_i(r)$ find new $FS_i'(r)$ and compute its amount of nectar;
 - Put the better one in $\{ FS_i'(r), FS_i(r) \}$ as a new position of the employed bees;
 }
 (5) For $j=1:K$ do (perform further local search by onlooker bees which take help from employed bees)
 {
 - A food source $FS_i(r)$ selected from $\{ FS_i(r) \}$ for every onlooker bee by apply roulette wheel selection method which pass the probability of choosing calculated from eq. 5;
 - A new food source $FS_i'(r)$ in the neighborhood of $FS_i(r)$ is find then calculate its nectar amount;
 - Consider the better one in $\{ FS_i'(r), FS_i(r) \}$ as a new position of the corresponding bees;
 }
 (6) A new food sources exploiting (scout bees randomly perform global search).
 For $i=1:K$ do (food sources)
 {
 If $FS_i(r) = \{ FS_i(r-1) \}$ then $CC_i = CC_i + 1$;
 If $CC_i = \text{Limit}$ (Limit= $K * N$) then
 {
 - The $FS_i(r)$ is abandon and the associated employed bee become scout;
 - Generate randomly a new $FS_i(r)$ and the scout becomes an employed bee again ;
 - $CC_i = 0$;
 }
 }
 - Register the best food source FS_{best} found so far whose gross-profit maximized and the sum of the total cost not exceeds the capacity of knapsack (suppose C calculated from multiply 0.25 by sum of the total cost); $r=r+1$;
 }
 7- Output: if the predefined end condition (corresponding to access maximum of iteration) are met then return FS_{best} .

5- Experimental results

The proposed algorithm was implemented by matlab 8.3 that contribute in building efficient investment plan with zero-one knapsack problem. Initialize the parameters of ABC algorithm which illustrates in table (1).

Table (1) parameters of ABC algorithm

| Parameter | Value |
|-------------------------|-------|
| Iteration maximum | 400 |
| Number of employed bees | 100 |
| Number of onlooker bees | 100 |
| Number of scouts bees | 10 |

There are n projects each one has profit and cost for explain the steps of proposed algorithm in simplicity get the example with eight projects, in table (2) lists the sales-revenue of the projects.

Table (2) the sales-revenue for each project

| Project number | the sales-revenue of the projects |
|----------------|-----------------------------------|
| project1 | 92045 |
| project 2 | 90789 |
| project 3 | 58047 |
| project 4 | 106336 |
| project 5 | 49013 |
| project 6 | 50292 |
| project 7 | 51336 |
| project 8 | 105390 |

Table (3) illustrate cost-making which consist from four types: cost of direct materials, the cost of direct wages, variable industrial costs and non-variable industrial costs for each project.

Table (3) cost-making for the projects

| No. project | Project 1 | Project 2 | Project 3 | Project 4 | Project 5 | Project 6 | Project 7 | Project 8 |
|-------------------------------|-----------|-----------|-----------|-----------|-----------|-----------|-----------|-----------|
| Cost –making type | | | | | | | | |
| cost of direct materials | 5475 | 5353 | 6096 | 7496 | 7866 | 7503 | 5251 | 6832 |
| cost of direct wages | 5606 | 6765 | 7775 | 6211 | 7815 | 7220 | 5198 | 5431 |
| variable industrial costs | 7600 | 5967 | 7289 | 7642 | 6516 | 6412 | 7055 | 5885 |
| non-variable industrial costs | 6965 | 7720 | 7409 | 7155 | 7082 | 6800 | 6933 | 6134 |

Then calculate the total cost from summation of cost-making for each project, the result explained in table(4).

Table (4) summation of cost-making for each project

| No. project | Project1 | Project2 | Project3 | Project4 | Project5 | Project6 | Project7 | Project8 |
|-------------|----------|----------|----------|----------|----------|----------|----------|----------|
| Total cost | 25646 | 25805 | 28569 | 28504 | 29279 | 27935 | 24437 | 24282 |

Then compute gross profit by subtract total cost from sale-revenue showed in table(5).

Table (5) total profit for the projects

| No. project | Project1 | Project2 | Project3 | Project4 | Project5 | Project6 | Project7 | Project8 |
|-------------------|----------|----------|----------|----------|----------|----------|----------|----------|
| Total profit-loss | 66399 | 64984 | 29478 | 77831 | 19734 | 22357 | 29899 | 81108 |

In this paper the initial solutions are not generated at total cost for each project to find the relation, the randomly, obtained from eq.3 by divide gross-profit result in table (6).

Table (6) the relation between gross profit to the total cost

| No. project | Project1 | Project2 | Project3 | Project4 | Project5 | Project6 | Project7 | Project8 |
|--------------|----------|----------|----------|----------|----------|----------|----------|----------|
| Relation (R) | 2.5891 | 2.5183 | 1.0318 | 2.7305 | 0.6740 | 0.8003 | 1.1007 | 3.3403 |

then calculate the probability of each project from eq.4 illustrated in table (7).

Table (7) the probability of presence each project in initial solution

| No. project | Project1 | Project2 | Project3 | Project4 | Project5 | Project6 | Project7 | Project8 |
|-------------|----------|----------|----------|----------|----------|----------|----------|----------|
| Probability | 0.3502 | 0.3407 | 0.1396 | 0.3694 | 0.0912 | 0.1083 | 0.1489 | 0.4518 |

Generate the initial solution depending on the probabilities computed in previous step, then calculate gross -profit-loss and total cost for each solutions explained in table (8).

Table (8) the initial solution

| initial solution | cost | profit |
|-------------------|--------|--------|
| [0,0,1,1,0,1,1,0] | 109445 | 156565 |
| [1,0,1,0,1,0,1,1] | 132213 | 223618 |
| [1,1,0,1,1,1,0,0] | 137169 | 251305 |
| [1,1,0,0,1,1,1,1] | 157384 | 281481 |
| [1,0,0,0,0,1,0,0] | 53581 | 88756 |
| [0,0,0,0,0,0,1,0] | 24437 | 26899 |
| [1,1,0,1,1,0,1,0] | 133671 | 255847 |
| [0,1,0,0,0,0,0,0] | 25805 | 64984 |
| [0,0,0,0,0,0,1,1] | 48719 | 108007 |
| [0,1,1,0,0,0,0,0] | 54374 | 94462 |

Take the solutions which corresponding to the total cost not exceeds the capacity of knapsack obtained from multiply 0.25 by the sum of the total cost, and save as the best solution, then compute the value of fitness from eq.5. for the generated solution from divide the gross profit for each string at the sum of gross profit then passing the value of selection probability computed from eq.5 to the roulette wheel selection, then select new food position and evaluate this position by compute the value of nectar amount and compare with the previous food positions, then change with the worst and continue with the rest steps of ABC algorithm until reach to the stopping criteria of algorithm. In table (9) illustrates the final solution by select subset of best project whose maximum

gross - profit and sum of the total cost not exceed capacity of knapsack.

Table (9) the best solution from apply ABC algorithm

| Best solution | Cost | Profit | Executing time |
|-------------------|-------|--------|----------------|
| [1,0,0,0,0,1,0,0] | 53581 | 136025 | 0.159606 |

To compare ABC algorithm with other method such as genetic algorithm which its parameters explained in table (10).

Table (10) parameter of GA

| Parameter | Value |
|--------------------------|--------------------------|
| Itermax | 400 |
| Number of chromosome | 10 |
| Type of selection | Roulette wheel selection |
| Type of crossover | One point |
| Type of mutation | Flip bit |
| Probability of crossover | 0.6 |
| Probability of mutation | 0.1 |

The result from comparison two methods with examples with 8, 10, 25, 50, and 100 projects for finding the best solution showed in table (11), which explains the efficiency of ABC algorithm in access to the good quality solution than applying GA.

Table (11) comparison between applying ABC algorithm and genetic algorithm

| No. of project | Capacity of knapsack | Sum of gross-profit | | Sum of the total cost | | Executing time | |
|----------------|----------------------|---------------------|--------|-----------------------|--------|----------------|---------------|
| | | ABC | GA | ABC | GA | ABC | GA |
| 8 | 53614.75 | 136025 | 100842 | 53581 | 53561 | 0.159606 sec | 3.046434 sec |
| 10 | 64770.75 | 36169 | 30747 | 58702 | 58191 | 0.435974 sec | 40.988467 sec |
| 25 | 1647715.5 | 260177 | 256249 | 163993 | 163453 | 0.598020 sec | 69.291353 sec |
| 50 | 325416.75 | 685087 | 577105 | 324473 | 325197 | 1.878421 sec | 95.296004 sec |
| 100 | 645463 | 1345983 | 988996 | 645450 | 645462 | 16.623129 sec | 186.56857 sec |

6- Conclusion

In this paper a proposed ABC algorithm used with knapsack problem to select subset of project with maximum gross-profit and the sum of total cost for them not exceeds the capacity of knapsack. The

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obtained result shows the highly speed ,accuracy, and good quality of solution from applying ABC algorithm compared with genetic algorithm to find best solution that contribute in put optimal investment plan.

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انتقاء المشاريع في مسألة حقيبة نابساك باستخدام خوارزمية مستعمرة النحل الاصطناعية

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الملخص

تعد مسألة حقيبة نابساك واحدة من مسائل الأمثلية التي تهدف الى اختيار العناصر التي تمتلك الفائدة الاكبر ووزنها لا يعبر حد الحقيبة. تُقدم في هذا البحث خوارزمية النحلة لاختيار مجموعة جزئية من المشاريع وتمثيلها بمسألة الحقيبة لوضع الخطة الاستثمارية الأفضل التي تحقق أعلى الفوائد ضمن الكلف المحددة، وهذه الخطة هي إحدى تطبيقات المجال المالي. وقد بيّنت النتائج المستحصلة من الخوارزمية المقترحة والممثلة باستعمال الماتلاب إمكانية إيجاد الحل الأفضل بدقة وسرعة عالية مقارنة بالخوارزمية الجينية.